

State Space Reductions

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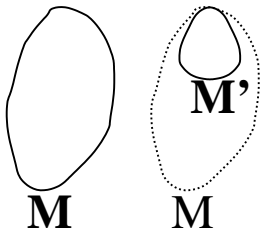
State Space Reductions project group
Parallel and Distributed Systems Laboratory

Seminar of ParaDiSe
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Contents

- efficient checking of p.o.r. properties for LTL model checking
- LTL model checking based on regular expressions
- distributed algorithms for SCCs decomposition

Partial Order Reduction for LTL



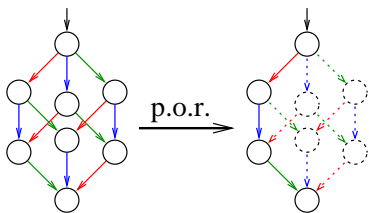
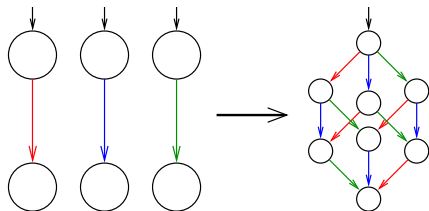
- $\varphi \in \text{LTL-X}: M \models \varphi \Leftrightarrow M' \models \varphi$
- construction of M' :
 - start from the initial state
 - for each reached state, explore a subset of enabled transitions

Partial Order Reduction for LTL

Conditions imposed on ample sets:

- C0: $ample(s) = \emptyset \Rightarrow enabled(s) = \emptyset$
- C1: every path in the full state space M starting at s can not have transition dependent on some transition in $ample(s)$ without a transition in $ample(s)$ occurring first
- C2: $ample(s)$ has visible transition $\Rightarrow ample(s) = enabled(s)$
- C3: for each cycle in M' holds: it contains state where α is enabled $\Rightarrow \alpha$ is included in ample set of some state in cycle

Partial Order Reduction for LTL



Partial Order Reduction for LTL – Improvement

- in general approach, only transitions of a single process are taken as a candidate to an ample set
- if more dependencies among transitions occur, p.o.r. achieves small reduction
- an approach assuming ample sets contained transitions from different processes proposed

LTL Based on Regular Expressions

My goal: to propose an approach to LTL model checking which:

- doesn't rely on construction of Büchi automaton
- isn't forced to store all states
- is exploration independent

LTL Based on Regular Expressions

FOMLO = First Order Monadic Logic.

- formulae express properties of infinite words over finite alphabet
- example of a FOMLO formula:
$$f = \forall x_1 : \exists x_2. x_2 > x_1 : x_2 \in P_a$$
- LTL and FOMLO are expressively equivalent (GFa is equivalent to f)

LTL Based on Regular Expressions

SFRE = Star-free Regular Expressions.

- SFRE expressions are built from letters of an alphabet, \emptyset and operands \cdot , $+$ and \neg
- Thomas (1979) has proven that FOMLO (LTL) and SFRE are expressively equivalent:
 - thus, for each FOMLO f there is an equivalent expression of the form $\bigcup U_i \cdot V_i^\omega$, where U_i, V_i are SFRE
 - equivalently, for each LTL formula φ , there is an equivalent expression of the form $E_\varphi = \bigcup U_i \cdot V_i^\omega$

LTL Based on Regular Expressions

Proposed approach: for a given model M and given LTL formula φ :

- find expression $E_{\neg\varphi}$
- test whether a behaviour of M satisfies $E_{\neg\varphi}$:
 - if so, $M \not\models \varphi$ and such a behaviour is a counterexample
 - if not so, $M \models \varphi$

Distributed SCCs Decomposition

Usage in model checking:

- LTL (property specified by GBA, Müller, or Street automaton)
- τ -confluence reduction
- probabilistic model checking (???)

Distributed SCCs Decomposition

Existing algorithms:

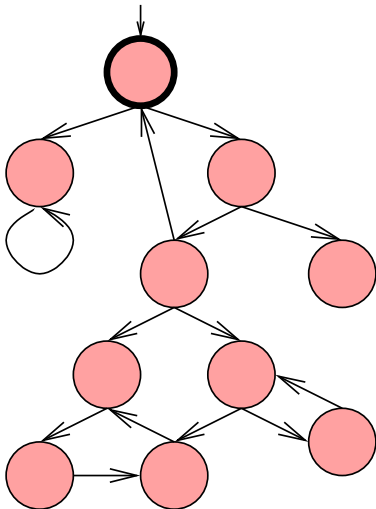
- sequential (Tarjan) – efficient distribution unknown
- FWD+BWD traversal
- colouring & colour prioritising (S. Orzan)

Distributed SCCs Decomposition

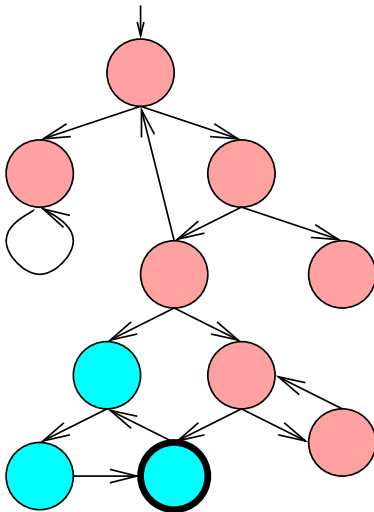
Our (my + Jiřík's) proposals:

- based on colouring a graph from all vertices in sequence
- storing additional information to reduce the number of reachabilities to be performed
- usage of back-level edges/states
- variants with back-propagations

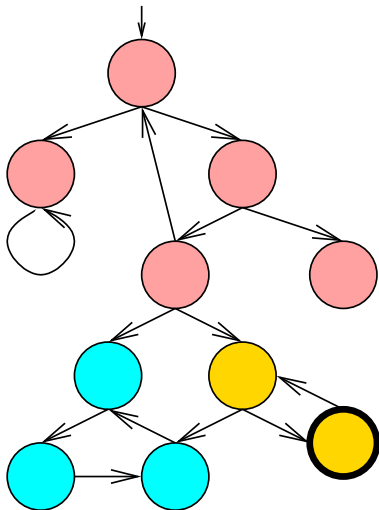
Algorithm Essence



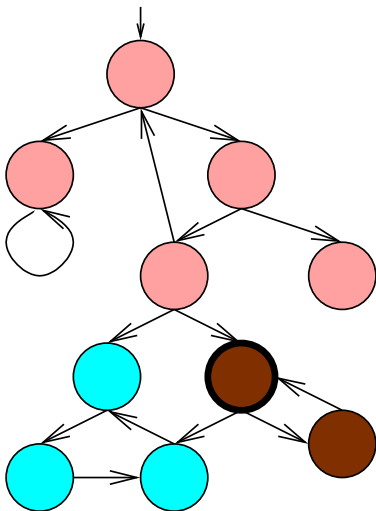
Algorithm Essence



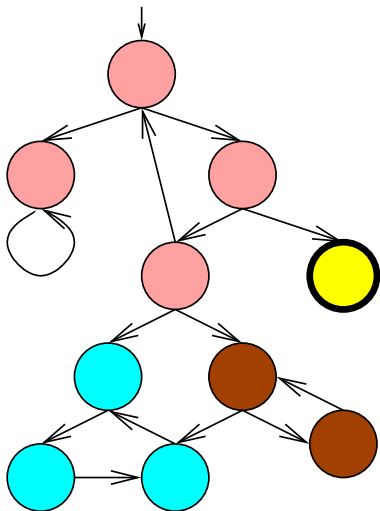
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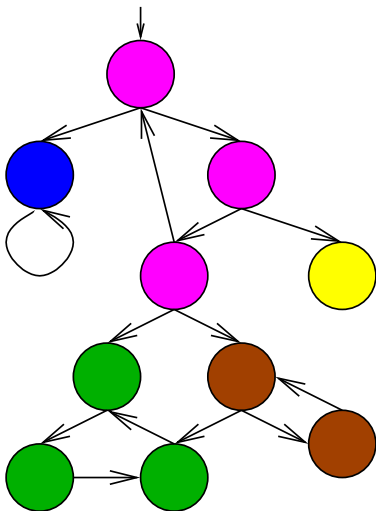
Algorithm Essence



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Algorithm Essence

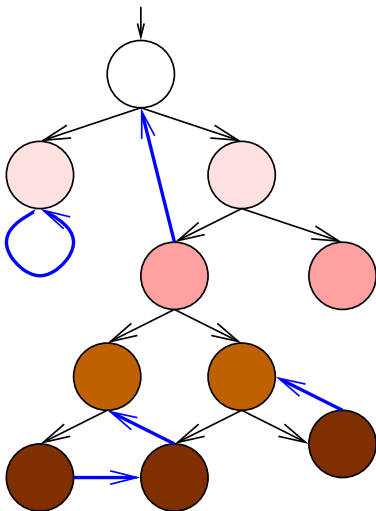


Distributed SCCs Decomposition

Our (Me + Jiřík) proposals:

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Back Level Edges



Distributed SCCs Decomposition

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Distributed SCCs Decomposition

Results:

- algorithms based on forward traversal only:
generally no great success
- improvement of FWD+BWD algorithm:
significant speedup
- research still in progress