## **Platform Dependent Verification** Engineering Verification Tools for 21st Century

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## Motivation



#### Complexity of developed systems is growing

- Numerical computation & text processing
- Database systems, computer games
- Simulations of physical systems
- Virtual reality

• ...



#### Parallelism and communication are essential



#### Correctness and reliability are critical measures.

## Flaws of concurrent systems

#### Concurrent systems are everywhere

- Network applications, data communication protocols
- Client-server systems
- Multi-threaded code
- ...

#### **Errors Related to Concurrency**

- Deadlocks
- Livelocks
- Underspecification
- Overspecification
- Assumption about execution speed



#### **Verification Methods**

• Simulations, Prototyping, Testing, Formal Verification

#### Formal Verification Benefits

- Early integration in design process.
- More effective (higher coverage).
- Reduction in verification time.
- Reduction in development costs and time-to-market.

## **Press-the-button!**

... a brief look at model-checking

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## Model Checking



#### ACM Turing Award 2007

- Edmund M. Clarke jr. (CMU, USA)
- Allen E. Emerson (Texas at Austin, USA)
- Joseph Sifakis (IMAG Grenoble, France)

#### Jury justification

"For their roles in developing Model-Checking into a highly effective verification technology, widely adopted in the hardware and software industries."

## Model Checking and State Space Explosion Problem



## Model Checking and State Space Explosion Problem



### PDMC Workshop Series



## Fight the state explosion problem by squeezing all the power out of the contemporary hardware

## What is needed ...

#### **Contemporary computing platforms**

- Distributed memory systems
- Shared memory parallelism
- GP GPU many-cores
- External memory
- ... and combinations.

#### **Efficient utilization requires**

- Parallel algorithms
- Algorithm engineering



# The DiVinE Story!

... platform dependent model checking with DiVinE

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## The DiVinE Story: DiVinE

#### What is DiVinE?

- Explicit-state automata-based LTL model checker.
- Can exploit aggregate computing power of multiple network-interconnected multi-cored workstations.

#### Where is it?

http://divine.fi.muni.cz



# The DiVinE Story!

... parallel algorithms

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#### Automata-based LTL Model Checking

• Accepting cycle detection in a directed graph.

#### **Optimal Serial Algorithms**

- Rely on **difficult-to-parallelize** DFS-postorder.
- Inefficient on parallel/distributed HW architectures and non-flat memory hierarchies.

#### Parallel scalable algorithms needed!

## The DiVinE Story: Parallel Scalable Algorithms

	Complexity	Optimality	On-The-Fly
Nested DFS	O(N+M)	Yes	Yes
OWCTY Algorithm			
general LTL properties	O(N.(N+M))	No	No
weak LTL properties	O(N+M)	Yes	No
MAP Algorithm	O(N.N.(N+M))	No	Yes
MAP-OWCTY Algorithm			
general LTL properties	O(N.(N+M))	No	Yes
weak LTL properties	O(N+M)	Yes	Yes
BLEDGE Algorithm	O(N.N.(N+M))	No	Yes
NEGC Algorithm	O(N.N.(N+M))	No	Yes

- N number of states
- M number of transitions

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#### Observations

- Distributed memory processing should not stay isolated.
- It must combine with other techniques.

#### Partial Order Reduction (POR)

• Key technique to fight state space explosion in explicit state model checking.

#### POR Principle

Orders execution of independent actions among processes of the distributed system.

Distributed system definition





#### POR Principle

Orders execution of independent actions among processes of the distributed system.



#### **POR Principle**

Suppose actions a1, b1 and b2 are independent.



#### **POR Principle**

By postponing execution of selected independent action (e.g. a1) the state space graph is reduced.



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#### **POR Principle**

Infinite postponing is avoided if every cycle contains at least one fully expanded state  $\Rightarrow$  cycle proviso.



#### **POR Principle**

To detect the cycle, **depth-first search stack** is used in the standard sequential approach.



Suppose a directed graph.



First, indegree is computed for every vertex.



Vertices with zero indegree are "eliminated".



Again, vertices with zero indegree are "eliminated".



Vertices whose indigree has decreased, are marked and removed.

















All cycles are covered.



Generation starts from an initial state.

0

Initial part of the state space is generated using action-postponing principle.



Topological sort proviso is applied to mark at least one state on every cycle.



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Marked states are fully re-expanded.



New states are discovered outside the previously generated part of the state space.



New part of the reduced state space is generated using action-postponing principle.



Topological sort proviso is applied to mark at least one state on every cycle.



Marked states are fully re-expanded.



New state is discovered outside the previously generated parts of the state space.



Additional part of the reduced state space is generated using action-postponing principle.



Topological sort proviso is applied to mark at least one state on every cycle.



Marked states are fully re-expanded.



No new states are generated, hence the reduced state space has been fully constructed.



# The DiVinE Story!

... algorithm engineering

#### Algorithm engineering

"Efforts must be made to ensure that promising algorithms discovered by the theory community are implemented, tested and refined to the point where they can be usefully applied in practice."

[Aho et al. [1997], Emerging Opportunities for Theoretical Computer Science]

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#### In other words ...

... implementation matters.



## The DiVinE Story: Message Aggregation

#### Principle

Individual messages



are sorted according the target and combined





#### Buffers are flushed when

- explicitly requested by underlying algorithm, or
- maximal number of messages in buffer is reached, or
- node is otherwise idle (empty queue), or
- messages in the buffer are too old.

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## The DiVinE Story: Flushing Buffers





#### Graph Exploration Loop

- States to be explored locally are taken from a local queue.
- Queue contains locally generated states and states received from network.

#### When to process incoming messages?



#### **Graph Exploration Loop**

- States to be explored locally are taken from a local queue.
- Queue contains locally generated states and states received from network.

#### When to process incoming messages?

- High rate  $\rightarrow$  CPU load.
- $\bullet$  Low rate  $\rightarrow$  increased memory demands.

## The DiVinE Story: Performance of OWCTY Algorithm



In cooperation with VU Amsterdam.

Cores	Runtime (sec)	Efficiency
1	631.7	100%
64	13.3	74%
128	7.4	67%
256	5.0	49%

## The DiVinE Story: Engineering on Other Platforms

#### Threading in Shared-Memory

- Avoid false sharing.
- Efficient memory allocation and deallocation.
- Lock-free communication data structures.

#### Shared-Memory Memory Access

- Memory access optimized hash tables.
- Compact data representation.

#### External Memory (disks)

- Delayed duplicate detection.
- Recompute rather than communicate.

#### GP GPU Many Cores

- Avoid hash-based graph partitioning.
- LTL model checking as matrix-vector multiplication.

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## What's Next

... some conclusions

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#### **General observation**

- Gap between pseudo-code and implementation is widening.
- Implementations need to be tuned for individual platforms.
- We should learn to appreciate engineering solutions.

#### Explicit state verification future

- Combination with static analysis and symbolic approaches.
- Increased importance of techniques that trade space for time.
- Platform dependent state space reductions.

## The End

... thank you for your attention

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